

INSTRUCTOR'S SOLUTIONS MANUAL

DISCRETE MATHEMATICS FIFTH EDITION

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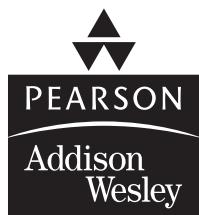
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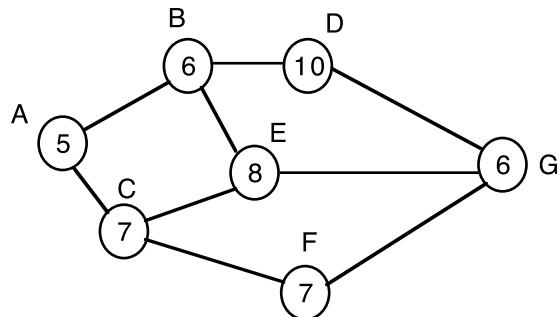
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Chapter 1

An Introduction to Combinatorial Problems and Techniques

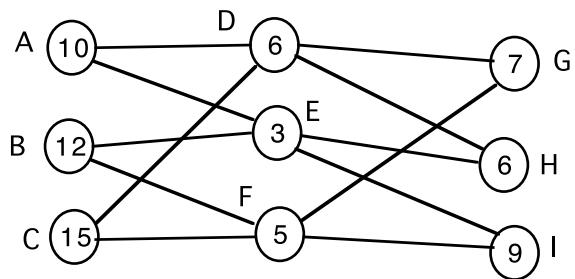
1.1 THE TIME TO COMPLETE A PROJECT

- 2. 31; A-B-E-G
- 4. 39; A-C-G-H
- 6. 16; B-D-F-H
- 8. 27; A-D-E-H
- 10. 27; A-B-D-G

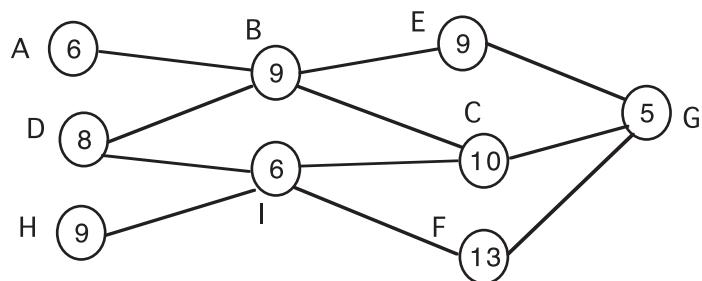


Chapter 1 An Introduction to Combinatorial Problems and Techniques

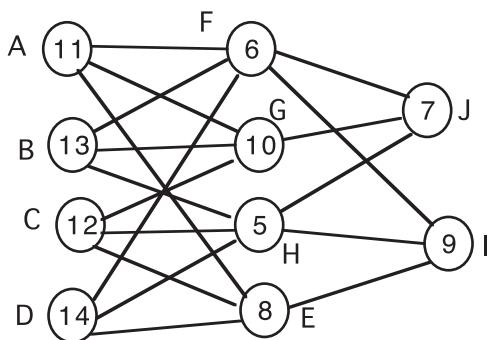
12. 29; C-F-I



14. 33; H-I-F-G



16. 31; D-E-I



18. 20 minutes

1.2 A Matching Problem

1.2 A MATCHING PROBLEM

- | | | | |
|-----------------------|--------------------------|-----------------|-----------------|
| 2. 720 | 4. 210 | 6. 84 | 8. 1680 |
| 10. 19,958,400 | 12. $\frac{1}{8}$ | 14. 5040 | 16. 126 |
| 18. 210 | 20. 119 | 22. 1320 | 24. 5040 |
| 26. 240 | 28. 1200 | | |

1.3 A KNAPSACK PROBLEM

- | | | | |
|--|-----------------|----------------|---------------|
| 2. T | 4. F | 6. F | 8. F |
| 10. T | 12. T | 14. T | 16. no |
| 18. yes; 32 | | | |
| 20. $\emptyset, \{1\}, \{2\}, \{3\}, \{4\}, \{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\}, \{1, 2, 3\}, \{1, 2, 4\}, \{1, 3, 4\}, \{2, 3, 4\}, \{1, 2, 3, 4\}; 16$ | | | |
| 22. 128 | 24. 1024 | 26. 256 | 28. 26 |
| 30. $\{1, 4, 6, 7, 8, 9, 10, 11, 12\}$ | | | |

1.4 ALGORITHMS AND THEIR EFFICIENCY

- | | |
|--|--------------------------------|
| 2. yes; 0 | 4. no |
| 6. no | 8. -1, 9, 84; 3, 17, 84 |
| 10. -4, -4, 41, 95; 2, 11, 33, 95 | 12. 111000 |
| 14. 001010 | |

Chapter 1 An Introduction to Combinatorial Problems and Techniques

16.	k	j	a_1	a_2	a_3	
	3		1	1	1	
	2		1	1	1	
	1		1	1	1	
	0		1	1	1	

18.	k	j	a_1	a_2	a_3	a_4	
	4		1	1	1	0	
			1	1	1	1	

20. The circled numbers in the table below indicated the items being compared.

	a_1	a_2	a_3	a_4	j	k
	23	5	(17)	(12)	1	3
	23	(5)	(12)	17		2
	(23)	(5)	12	17		1
	5	23	(12)	(17)	2	3
	5	(23)	(12)	17		2
	5	12	(23)	(17)	3	3
	5	12	17	23		

22. The circled numbers in the table below indicated the items being compared.

	a_1	a_2	a_3	a_4	a_5	j	k
	88	2	75	(10)	(48)	1	4
	88	2	(75)	(10)	48		3
	88	(2)	(10)	75	48		2
	(88)	(2)	10	75	48		1
	2	88	10	(75)	(48)	2	4
	2	88	(10)	(48)	75		3
	2	(88)	(10)	48	75		2
	2	10	88	(48)	(75)	3	4
	2	10	(88)	(48)	75		3
	2	10	48	(88)	(75)	4	4
	2	10	48	75	88		

24. 6.5 years, 2.7 seconds

26. 2.3×10^{10} years, 12.5 seconds

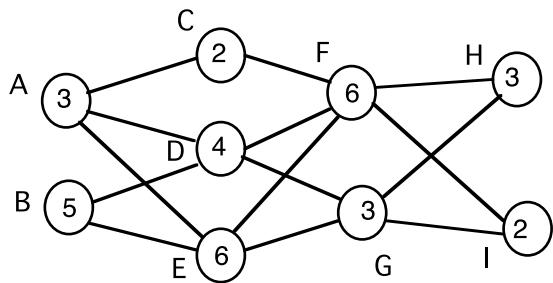
28. $4n - 3$

30. $3n - 2$

32. $-4, -4, 41, 95$

SUPPLEMENTARY EXERCISES

2. 20; B-E-F-H



4. 336

6. 40

8. 14040

10. T

12. F

14. T

16. T

18. 16

20. no

22. yes; 0

24. $-5, 7, 7, 88$ 26. $\emptyset, \{4\}, \{3\}, \{3, 4\}, \{2\}, \{2, 4\}, \{2, 3\}, \{2, 3, 4\}, \{1\}, \{1, 4\}, \{1, 3\}, \{1, 3, 4\}, \{1, 2\}, \{1, 2, 4\}, \{1, 2, 3\}, \{1, 2, 3, 4\}$ 28. 4.92×10^8 years

30. 4

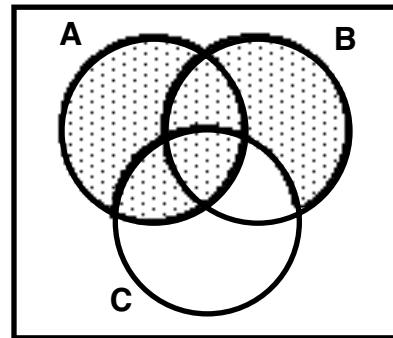
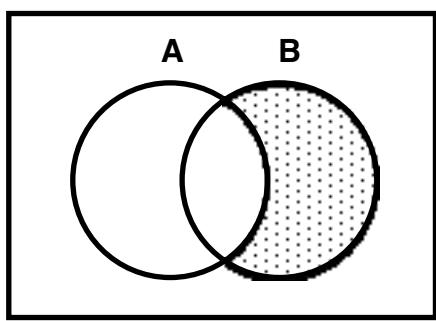
32. $4r - 3$

Chapter 2

Sets, Relations, and Functions

2.1 SET OPERATIONS

2. $A \cup B = \{1, 2, 4, 5, 6, 7, 9\}$, $A \cap B = \{1, 4, 6, 9\}$, $A - B = \emptyset$, $\overline{A} = \{2, 3, 5, 7, 8\}$, and $\overline{B} = \{3, 8\}$
4. $A \cup B = \{2, 3, 4, 5, 6, 7, 8, 9\}$, $A \cap B = \{7, 9\}$, $A - B = \{3, 4, 6, 8\}$, $\overline{A} = \{1, 2, 5\}$, and $\overline{B} = \{1, 3, 4, 6, 8\}$
6. $\{(3, 1), (3, 2), (3, 3), (4, 1), (4, 2), (4, 3), (5, 1), (5, 2), (5, 3)\}$
8. $\{(p, a), (p, c), (p, e), (q, a), (q, c), (q, e), (r, a), (r, c), (r, e), (s, a), (s, c), (s, e)\}$
- 10.
- 12.



14. $A = \{1, 2\}$, $B = \{1, 3\}$, and $C = \{1\}$
16. $A = \{1, 2\}$, $B = \{1, 3\}$, and $C = \{1\}$

18. A

20. \emptyset

22. $A \cap B$

24. $A \cup \overline{B}$

26. The equality $A - B = B - A$ holds if and only if $A = B$.

28. The equality $A \cap B = A$ holds if and only if $A \subseteq B$.

2.2 EQUIVALENCE RELATIONS

2. reflexive and symmetric

4. reflexive, symmetric, and transitive

6. reflexive and symmetric

8. none

10. reflexive and symmetric

12. reflexive and transitive

14. The equivalence class of R containing ABCD consists of the string ABC and the strings of 4 letters having A as their first letter and C as their third letter. There are $26^2 = 676$ distinct equivalence classes of R .

16. The equivalence class of R containing $\{1, 2, 3\}$ is the set containing the following four elements of S : $\{1, 3\}$, $\{1, 2, 3\}$, $\{1, 2, 3, 4\}$, and $\{1, 3, 4\}$. There are 8 different equivalence classes of R , namely the sets consisting of the elements S , $S \cup \{2\}$, $S \cup \{4\}$, and $S \cup \{2, 4\}$ for every $S \subseteq \{1, 3, 5\}$.

18. The equivalence class of R containing $(5, 8)$ is the set

$$\{(x_1, x_2) : \text{each } x_i \text{ is an integer and } x_1 - x_2 = 5 - 8\}.$$

There are infinitely many distinct equivalence classes of R , namely, the sets of the form

$$\{(x_1, x_2) : \text{each } x_i \text{ is an integer and } x_1 - x_2 = k\},$$

where $k = 0, \pm 1, \pm 2, \pm 3, \dots$

20. $\{(1, 1), (1, 3), (1, 6), (3, 1), (3, 3), (3, 6), (6, 1), (6, 3), (6, 6), (2, 2), (2, 5), (5, 2), (5, 5), (4, 4)\}$

24. The equivalence classes have the form $E_1 \times E_2$, where E_i is an equivalence class of R_i .

28. There are 5 partitions of a set with three elements.

32. Let S be a nonempty set and R an equivalence relation on S . Then there is a function f with domain S such that $s_1 R s_2$ if and only if $f(s_1) = f(s_2)$.

2.3 PARTIAL ORDERING RELATIONS

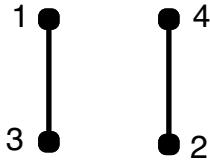
2. not antisymmetric

4. partial ordering

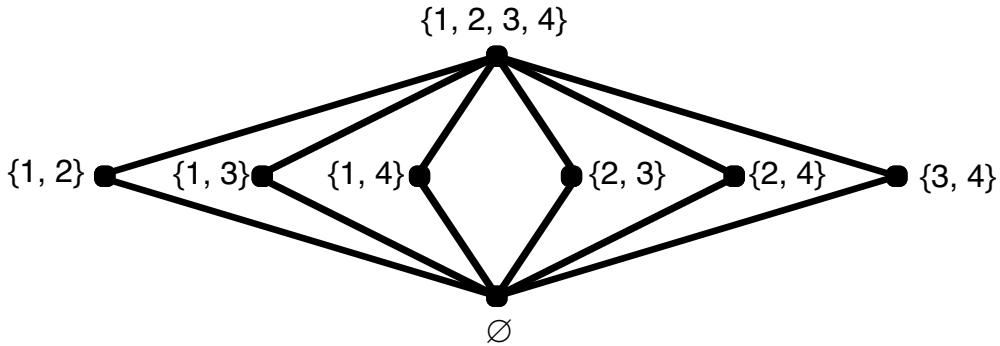
6. not antisymmetric

8. not antisymmetric

10.



12.



14. $R = \{(a, a), (b, b), (b, a), (c, c), (c, b), (c, a), (d, d), (d, a)\}$

16. $R = \{(1, 1), (2, 2), (2, 1), (2, 4), (4, 4), (8, 8), (8, 4)\}$

18. The maximal elements are $\{1\}$, $\{2\}$, and $\{3\}$; the only minimal element is $\{1, 2, 3\}$.

20. The only minimal element is 0; there are no maximal elements.

22. One possible sequence is 1, 3, 2, 4.

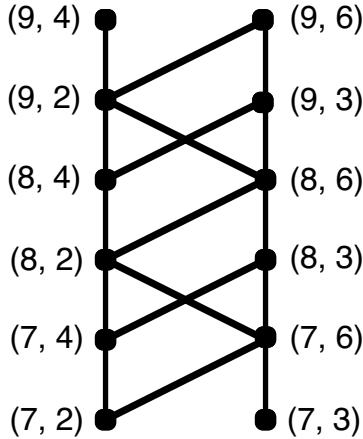
24. One possible sequence is 1, 3, 2, 6, 4, 12.

26. Let S denote the set of residents of Illinois and R be defined so that $x R y$ means that x is a sister of y .

28. Every prime integer is a minimal element; there are no maximal elements.

30. $(1, 1), (1, 2), (1, 3), (1, 4), (2, 1), (2, 2), (2, 3), (2, 4), (3, 1), (3, 2), (3, 3), (3, 4)$

32.



36. (a) Suppose that y is element of S such that $x R y$ is false. If there is no element y_1 in S such that $y_1 R y$, then y is a minimal element of S , contradicting that x is the unique minimal element of S . Thus there must be such an element y_1 . If there is no element y_2 in S such that $y_2 R y_1$, then y_1 is a minimal element of S , another contradiction. So there must be such an element y_2 . Because S is finite, continuing in this manner must produce a minimal element y_k of S different from x . Because x is the unique minimal element of S , the assumption that there is an element y of S such that $x R y$ is false must be incorrect. Thus $x R s$ is true for every $s \in S$.
- (b) Let Z denote the set of integers and $S = Z \cup \{\emptyset\}$. Let R be the relation defined on S by $x R y$ if and only if one of the following holds: (i) $x, y \in Z$ and $x \leq y$, or (ii) $x = y = \emptyset$. Then \emptyset is the unique minimal element in S , but $\emptyset R z$ is false for every $z \in Z$.

40. 2^n 42. $2^n \cdot 3^{n(n-1)/2}$